

# Graphs in $-1$ colors

## Learning to Live with Stanley's Acyclicity Theorem

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# The Main Result

Recall that a coloring of a graph  $G$  is a function  $\kappa$  which assigns every vertex in  $G$  a natural number. A coloring is proper if any two vertices which are connected by an edge have distinct colors.

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## Theorem (Stanley's Acyclicity Theorem)

*For any graph  $G$  with  $n$  vertices, the number of acyclic orientations of  $G$  is  $(-1)^n \chi_G(-1)$ , where  $\chi_G$  is the chromatic polynomial of  $G$ .*

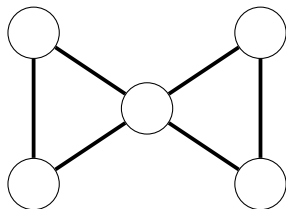
# Color by Numbers

To say that  $\kappa$  is a coloring of  $G$  in  $\lambda$  colors means that  $\kappa$  is a function  $\kappa : \{\text{Vertices of } G\} \rightarrow \{1, 2, \dots, \lambda\}$

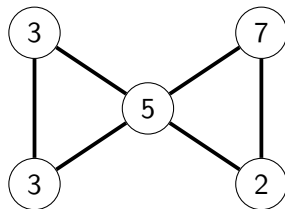
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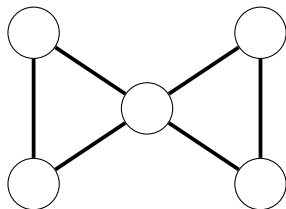
$G$  with coloring  $\kappa$



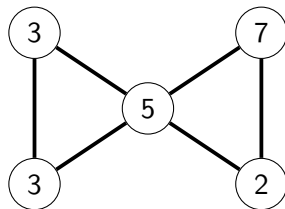
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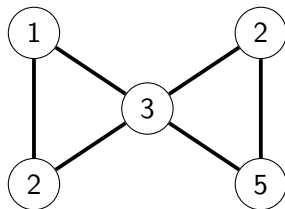
$G$  with coloring  $\kappa$



Note that  $\kappa$  is not a proper coloring since there are two vertices which are connected by an edge and have the same color.

## Coloring by Numbers (cont.)

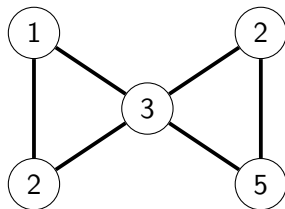
$G$  with proper coloring  $\kappa'$





## Coloring by Numbers (cont.)

$G$  with proper coloring  $\kappa'$



The chromatic polynomial for  $G$  is  $\chi_G(\lambda) = \lambda^5 - 6\lambda^4 + 13\lambda^3 - 12\lambda^2 + 4\lambda$ .

So for example, there are:

0 ways to properly color  $G$  in 2 colors

12 ways to properly color  $G$  in 3 colors

144 ways to properly color  $G$  in 4 colors

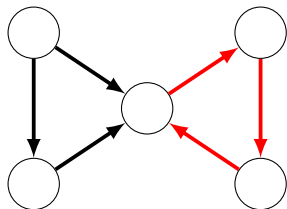
720 ways to properly color  $G$  in 5 colors

⋮

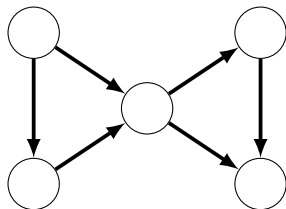
# What's an Acyclicity?

An orientation of a graph assigns a direction to each edge. The orientation is acyclic if it forms no cycles.

$G$  with cyclic orientation  $\sigma$



$G$  with acyclic orientation  $\sigma'$



As it turns out, there is an interesting relation between graph colorings and orientations.

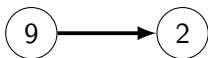
# A Connection Between Colors and Directions

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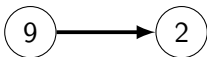
Specifically, each edge in  $G$  will point to the vertex whose coloring has a lower value.



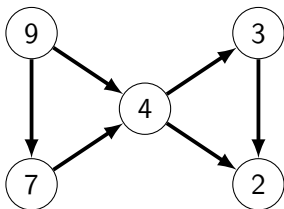
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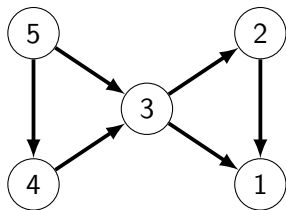
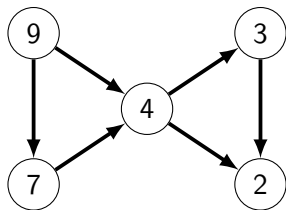


Any orientation created this way must be acyclic: the colors strictly decrease along each edge, so it would be impossible for a directed path to begin and end at the same vertex.



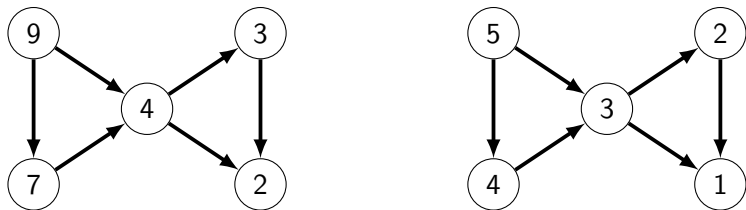
## A Connection Between Colors and Directions (cont.)

So, every proper  $\lambda$ -coloring of a graph  $G$  can be associated to an acyclic orientation of  $G$ . However, multiple proper colorings might be associated to the same orientation.



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Another way of phrasing this remark is that

$$\chi_G(\lambda) = \#\{(\kappa, \sigma) : \sigma \text{ is acyclic and } v_1 \xrightarrow{\sigma} v_2 \implies \kappa(v_1) > \kappa(v_2)\}$$

# What's Upsilon?

Let's consider a similar function,  $v_G$ , defined as

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Specifically, we can see that  $v_G(\mathbf{1})$  counts the number of acyclic orientations of  $G$ .

Our goal is to prove that  $(-1)^n v_G(\lambda) = \chi_G(-\lambda)$ , where  $n$  is the number of vertices in  $G$ .

# Contracting and Deleting

Recall that  $\chi$  is uniquely defined by the following conditions:

$$(1) \chi_o(\lambda) = \lambda$$

$$(2) \chi_{G \sqcup H}(\lambda) = \chi_G(\lambda) \cdot \chi_H(\lambda)$$

$$(3) \chi_G(\lambda) = \chi_{G-e}(\lambda) - \chi_{G/e}(\lambda)$$

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We will show that  $v$  has the following properties:

- (1')  $v_o(\lambda) = \lambda$
- (2')  $v_{G \sqcup H}(\lambda) = v_G(\lambda) \cdot v_H(\lambda)$
- (3')  $v_G(\lambda) = v_{G-e}(\lambda) + v_{G/e}(\lambda)$

## Contracting and Deleting (cont.)

$v$  satisfies (1') and (2') for the same reason that  $\chi$  does.

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First, if there are  $\lambda$  colors available, then there are  $\lambda$  ways to color a single vertex. For each coloring, there is only one way to “orient” a single vertex with no edges, so the total number of compatible pairs is  $\lambda$ .

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Second, if  $G$  and  $H$  are disjoint graphs, then any compatible pair for  $G$  put together with a compatible pair for  $H$  will make a compatible pair for  $G \sqcup H$ . Additionally, a compatible pair for  $G \sqcup H$  will remain a compatible pair when restricted to either  $G$  or  $H$ . So,  $v_{G \sqcup H}(\lambda) = v_G(\lambda) \cdot v_H(\lambda)$ .

## Contracting and Deleting (cont.)

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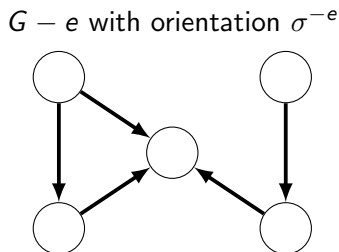
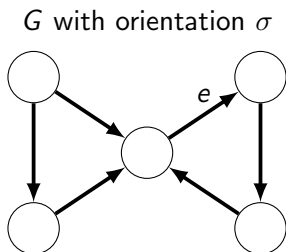
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Pick your favorite edge  $e$  in  $G$ . Consider  $\sigma^{-e}$  to be the orientation on  $G - e$  that agrees with  $\sigma$  on all edges other than  $e$  (it has no direction on  $e$  since  $e$  is not an edge in  $G - e$ .)

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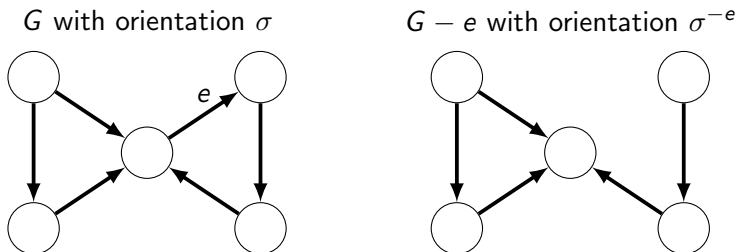
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We will first show that the mapping  $(\kappa, \sigma) \mapsto (\kappa, \sigma^{-e})$  between compatible pairs for  $G$  and compatible pairs for  $G - e$  is surjective for any choice of edge  $e$ .

# Proving Surjectivity

## Lemma

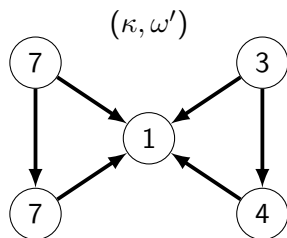
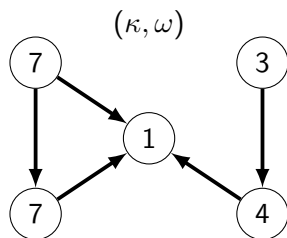
*For any compatible pair  $(\kappa, \omega)$  for  $G - e$ ,  $\omega$  can be extended to some  $\omega'$  on  $G$  such that  $(\kappa, \omega')$  is a compatible pair for  $G$ .*

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If  $u$  and  $v$  are the endpoints of  $e$  with  $\kappa(u) \neq \kappa(v)$ , then defining  $\omega'$  to point toward the vertex with the smaller color will result in  $(\kappa, \omega')$  being a compatible pair.

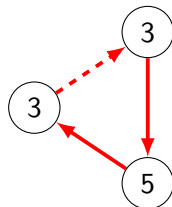
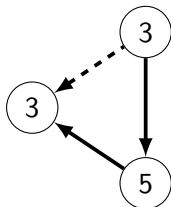


# Surjectivity, Subjectively

If  $\kappa(u) = \kappa(v)$ , then it is impossible for both choices of direction for  $e$  to result in a cycle, so we may pick whichever one makes  $\omega'$  acyclic.

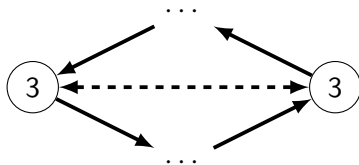
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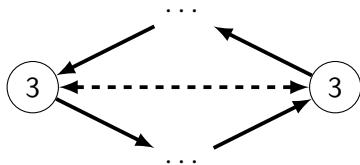
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The upshot of all of this is that the mapping  $(\kappa, \sigma) \mapsto (\kappa, \sigma^{-e})$  between compatible pairs for  $G$  and compatible pairs for  $G - e$  is surjective for any choice of edge  $e$ .

# Injectivity, Objectively

If this mapping is a bijection, then we are done.

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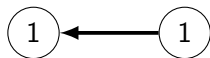
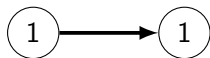
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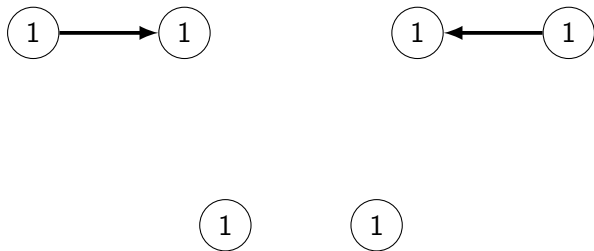
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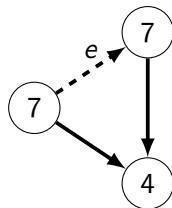
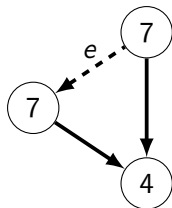
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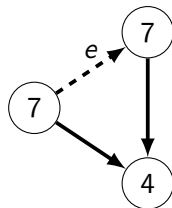
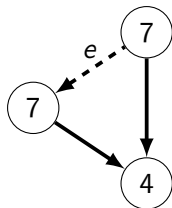
For two compatible pairs for  $G$  to be mapped to the same compatible pair for  $G - e$ , they must agree on every edge except for  $e$ . If  $e$  can point in either direction, then it must be that  $\kappa(u) = \kappa(v)$  since  $\kappa(u) \geq \kappa(v)$  and  $\kappa(u) \leq \kappa(v)$  (where  $u$  and  $v$  are the endpoints of  $e$ ).



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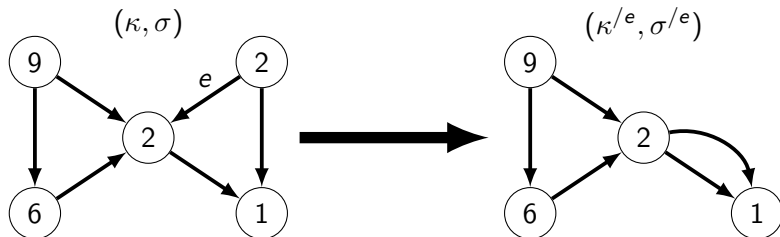


However, this means that a compatible pair is induced on  $G/e$ .



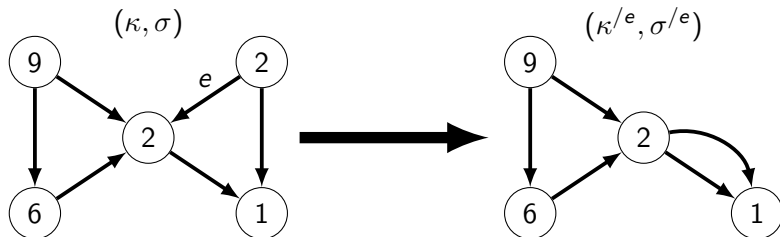
## Contractually Obligated Injectivity

If  $\kappa(u) = \kappa(v)$ , then we can easily define  $\kappa^{/e}$  and  $\sigma^{/e}$  on  $G/e$ . We demand that  $\kappa^{/e}$  and  $\sigma^{/e}$  agree with  $\kappa$  and  $\sigma$  wherever possible and  $\kappa^{/e}(w) = \kappa(u) = \kappa(v)$  where  $w$  is the vertex in  $G/e$  created by fusing  $u$  and  $v$  together.



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Therefore, any compatible pair for  $G/e$  can be extended to not just one compatible pair for  $G$ , but in fact two compatible pairs for  $G$ .

# Putting it All Together

Considering all of the compatible pairs for  $G$  (which by definition is  $v_G(\lambda)$ ), we've seen that they can all be produced by extending the  $v_{G-e}(\lambda)$  compatible pairs of  $G - e$ , and that  $v_{G/e}(\lambda)$  of those pairs admit either direction of  $e$  in the extension.

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So we may consider our map from before that maps compatible pairs of  $G$  to compatible pairs for  $G - e$ ,  $(\kappa, \sigma) \mapsto (\kappa, \sigma^{-e})$ . We have found that there are two pairs which are mapped to the same thing under this map for each compatible pair for  $G/e$  (of which there are  $v_{G/e}(\lambda)$ ) and moreover, these are the only pairs which are mapped to the same thing.

## Putting it All Together (cont.)

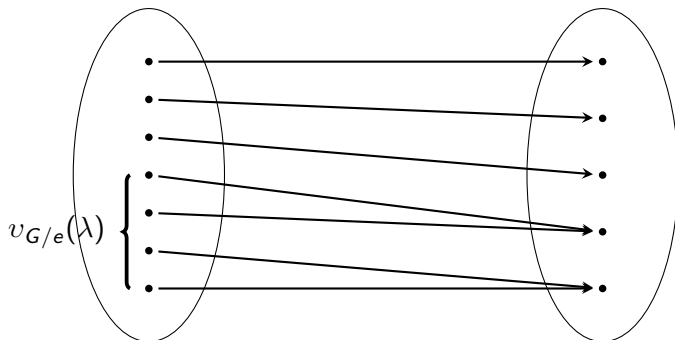
So, we can put some of the compatible pairs for  $G$  in 1-to-1 correspondence with the compatible pairs for  $G - e$  (of which there are  $v_{G-e}(\lambda)$ ) and have  $v_{G/e}(\lambda)$  compatible pairs left over.

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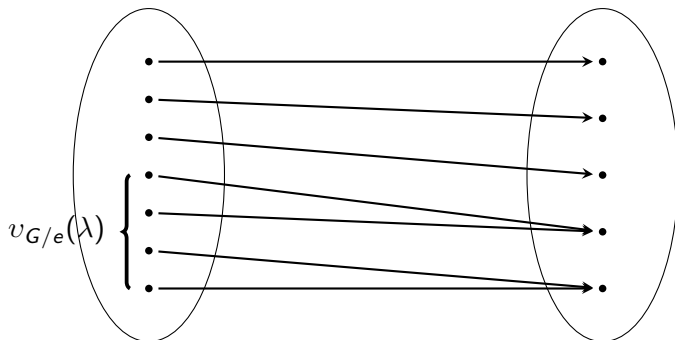


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Compatible pairs for  $G - e$



In other words, we have that  $v_G(\lambda) = v_{G-e}(\lambda) + v_{G/e}(\lambda)$ .

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Using this information, we can show that  $(-1)^n v_G(\lambda) = \chi_G(-\lambda)$ .



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$$(-1)^1 v_o(\lambda) = -\lambda \quad (1')$$

$$(-1)^{n+m} v_{G \sqcup H}(\lambda) = (-1)^n v_G(\lambda) \cdot (-1)^m v_H(\lambda) \quad (2')$$

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$$(-1)^{n+m} v_{G \sqcup H}(\lambda) = (-1)^n v_G(\lambda) \cdot (-1)^m v_H(\lambda) \quad (2')$$

And lastly,

$$\begin{aligned} (-1)^n v_G(\lambda) &= (-1)^n v_{G-e}(\lambda) + (-1)^n v_{G/e}(\lambda) \\ &= (-1)^n v_{G-e}(\lambda) - (-1)^{n-1} v_{G/e}(\lambda) \end{aligned} \quad (3')$$

## A Negative Number of Colors (cont.)

We know that  $\chi_G(-\lambda)$  is uniquely defined by the following properties

- (1)  $\chi_o(-\lambda) = -\lambda$ .
- (2)  $\chi_{G \sqcup H}(-\lambda) = \chi_G(-\lambda) \cdot \chi_H(-\lambda)$ .
- (3)  $\chi_G(-\lambda) = \chi_{G-e}(-\lambda) - \chi_{G/e}(-\lambda)$ .

We have shown that  $(-1)^n v_G(\lambda)$  satisfies each of these properties and so we have that

$$(-1)^n v_G(\lambda) = \chi_G(-\lambda)$$

Using this formula, we can note that by plugging in  $\lambda = 1$  we arrive at

$$(-1)^n v_G(1) = \chi_G(-1)$$

$v_G(1)$  is the number of acyclic orientations of  $G$  since there is only one way to color a graph in one color.

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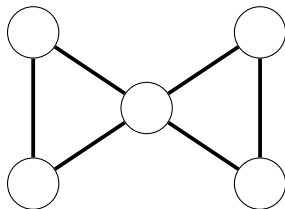
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And so in conclusion:

The number of acyclic orientations of  $G = v_G(1) = (-1)^n \chi(-1)$

## Example

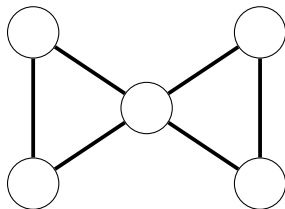
Consider our favorite graph  $G$ :



It has 6 edges and so it has  $2^6 = 64$  orientations. There are only two paths that can be cycles. If one of the two paths is a cycle, then there are  $2^3 = 8$  ways to orient the other half. There are two possible cycles and each cycle can be directed in one of two ways and so (taking into account the 4 ways in which there can be two cycles) we have that there are  $8 \cdot 2 \cdot 2 - 4 = 28$  orientations of  $G$  which are cyclic.

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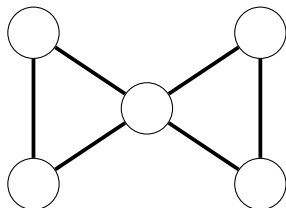


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Therefore  $G$  has  $64 - 28 = 36$  acyclic orientations.



## Example (cont.)



Recall that the chromatic polynomial for  $G$  is

$$\chi_G(\lambda) = \lambda^5 - 6\lambda^4 + 13\lambda^3 - 12\lambda^2 + 4\lambda.$$

$G$  has 5 vertices and so Stanley's Acyclicity Theorem tells us that  $G$  has  $(-1)^5 \chi_G(-1) = -(-1 - 6 - 13 - 12 - 4) = 36$  acyclic orientations.

Stanley, Richard P. “Acyclic Orientations of Graphs.” *Discrete Mathematics*, vol. 5, no. 2, 1973, pp. 171–178., doi:10.1016/0012-365x(73)90108-8.